

Key Management Interoperability Protocol Test Cases Version 1.3

Committee Note Draft 01 /
Public Review Draft 01

07 April 2016

Specification URIs

This version:

<http://docs.oasis-open.org/kmip/testcases/v1.3/cnprd01/kmip-testcases-v1.3-cnprd01.docx> (Authoritative)

<http://docs.oasis-open.org/kmip/testcases/v1.3/cnprd01/kmip-testcases-v1.3-cnprd01.html>

<http://docs.oasis-open.org/kmip/testcases/v1.3/cnprd01/kmip-testcases-v1.3-cnprd01.pdf>

Previous version:

N/A

Latest version:

<http://docs.oasis-open.org/kmip/testcases/v1.3/kmip-testcases-v1.3.docx> (Authoritative)

<http://docs.oasis-open.org/kmip/testcases/v1.3/kmip-testcases-v1.3.html>

<http://docs.oasis-open.org/kmip/testcases/v1.3/kmip-testcases-v1.3.pdf>

Technical Committee:

[OASIS Key Management Interoperability Protocol (KMIP) TC](https://www.oasis-open.org/committees/kmip/)

Chairs:

Tony Cox (tony.cox@cryptsoft.com), [Cryptsoft Pty Ltd.](http://cryptsoft.com/)

Saikat Saha (saikat.saha@oracle.com), [Oracle](http://www.oracle.com/)

Editors:

Tim Hudson (tjh@cryptsoft.com), [Cryptsoft Pty Ltd.](http://cryptsoft.com/)

Mark Joseph (mark@p6r.com), [P6R, Inc](http://www.p6r.com/)

Additional artifacts:

This document is one component of a Work Product that also includes:

* Test cases: <http://docs.oasis-open.org/kmip/testcases/v1.3/cnprd01/test-cases/kmip-v1.3/>

Related work:

This document replaces or supersedes:

* *Key Management Interoperability Protocol Test Cases Version 1.1*. Edited by Mathias Björkqvist and Tim Hudson. 27 January 2012. OASIS Committee Note 01. <http://docs.oasis-open.org/kmip/testcases/v1.1/kmip-testcases-v1.1.html>.
* *Key Management Interoperability Protocol Test Cases Version 1.2*. Edited by Tim Hudson and Faisal Faruqui. Latest version. <http://docs.oasis-open.org/kmip/testcases/v1.2/kmip-testcases-v1.2.html>.

This document is related to:

* *Key Management Interoperability Protocol Specification Version 1.3*. Edited by Kiran Thota and Tony Cox. Latest version: <http://docs.oasis-open.org/kmip/spec/v1.3/kmip-spec-v1.3.html>.
* *Key Management Interoperability Protocol Profiles Version 1.3.* Edited by Tim Hudson and Robert Lockhart. Latest version: <http://docs.oasis-open.org/kmip/profiles/v1.3/kmip-profiles-v1.3.html>.
* *Key Management Interoperability Protocol Usage Guide Version 1.3*. Edited by Judith Furlong. Latest version: <http://docs.oasis-open.org/kmip/ug/v1.3/kmip-ug-v1.3.html>.

Abstract:

This document is intended for developers and architects who wish to design systems and applications that interoperate using the Key Management Interoperability Protocol specification.

Status:

This document was last revised or approved by the OASIS Key Management Interoperability Protocol (KMIP) TC on the above date. The level of approval is also listed above. Check the “Latest version” location noted above for possible later revisions of this document.

Technical Committee (TC) members should send comments on this document to the TC’s email list. Others should send comments to the TC’s public comment list, after subscribing to it by following the instructions at the “[Send A Comment](https://www.oasis-open.org/committees/comments/index.php?wg_abbrev=kmip)” button on the TC’s web page at <https://www.oasis-open.org/committees/kmip/>.

Citation format:

When referencing this document the following citation format should be used:

**[kmip-testcases-v1.3]**

*Key Management Interoperability Protocol Test Cases Version 1.3.* Edited by Tim Hudson and Mark Joseph. 07 April 2016. OASIS Committee Note Draft 01 / Public Review Draft 01. <http://docs.oasis-open.org/kmip/testcases/v1.3/cnprd01/kmip-testcases-v1.3-cnprd01.html>. Latest version: <http://docs.oasis-open.org/kmip/testcases/v1.3/kmip-testcases-v1.3.html>.

Copyright © OASIS Open 2016. All Rights Reserved.

All capitalized terms in the following text have the meanings assigned to them in the OASIS Intellectual Property Rights Policy (the "OASIS IPR Policy"). The full [Policy](https://www.oasis-open.org/policies-guidelines/ipr) may be found at the OASIS website.

This document and translations of it may be copied and furnished to others, and derivative works that comment on or otherwise explain it or assist in its implementation may be prepared, copied, published, and distributed, in whole or in part, without restriction of any kind, provided that the above copyright notice and this section are included on all such copies and derivative works. However, this document itself may not be modified in any way, including by removing the copyright notice or references to OASIS, except as needed for the purpose of developing any document or deliverable produced by an OASIS Technical Committee (in which case the rules applicable to copyrights, as set forth in the OASIS IPR Policy, must be followed) or as required to translate it into languages other than English.

The limited permissions granted above are perpetual and will not be revoked by OASIS or its successors or assigns.

This document and the information contained herein is provided on an "AS IS" basis and OASIS DISCLAIMS ALL WARRANTIES, EXPRESS OR IMPLIED, INCLUDING BUT NOT LIMITED TO ANY WARRANTY THAT THE USE OF THE INFORMATION HEREIN WILL NOT INFRINGE ANY OWNERSHIP RIGHTS OR ANY IMPLIED WARRANTIES OF MERCHANTABILITY OR FITNESS FOR A PARTICULAR PURPOSE.

Table of Contents

[1 Introduction 6](#_Toc449973229)

[1.1 References (non-normative) 6](#_Toc449973230)

[2 KMIP Test Cases 7](#_Toc449973231)

[2.1 KMIP 1.3 Test Cases 7](#_Toc449973232)

[2.1.1 TC-CREG-2-13 7](#_Toc449973233)

[2.1.2 TC-OFFSET-1-13 8](#_Toc449973234)

[2.1.3 TC-OFFSET-2-13 8](#_Toc449973235)

[2.1.4 TC-OTP-1-13 8](#_Toc449973236)

[2.1.5 TC-OTP-2-13 8](#_Toc449973237)

[2.1.6 TC-OTP-3-13 8](#_Toc449973238)

[2.1.7 TC-OTP-4-13 9](#_Toc449973239)

[2.1.8 TC-OTP-5-13 9](#_Toc449973240)

[2.1.9 TC-Q-CAP-1-13 9](#_Toc449973241)

[2.1.10 TC-Q-CAP-2-13 9](#_Toc449973242)

[2.1.11 TC-Q-CREG-1-13 9](#_Toc449973243)

[2.1.12 TC-Q-PROF-1-13 9](#_Toc449973244)

[2.1.13 TC-Q-PROF-2-13 9](#_Toc449973245)

[2.1.14 TC-Q-PROF-3-13 10](#_Toc449973246)

[2.1.15 TC-Q-RNGS-1-13 10](#_Toc449973247)

[2.1.16 TC-Q-RNGS-2-13 10](#_Toc449973248)

[2.1.17 TC-Q-RNGS-3-13 10](#_Toc449973249)

[2.1.18 TC-Q-RNGS-4-13 10](#_Toc449973250)

[2.1.19 TC-Q-RNGS-5-13 10](#_Toc449973251)

[2.1.20 TC-Q-RNGS-6-13 10](#_Toc449973252)

[2.1.21 TC-Q-S2C-1-13 11](#_Toc449973253)

[2.1.22 TC-Q-S2C-2-13 11](#_Toc449973254)

[2.1.23 TC-Q-S2C-PROF-1-13 11](#_Toc449973255)

[2.1.24 TC-Q-S2C-PROF-2-13 11](#_Toc449973256)

[2.1.25 TC-Q-VAL-1-13 11](#_Toc449973257)

[2.1.26 TC-Q-VAL-2-13 11](#_Toc449973258)

[2.1.27 TC-RNG-ATTR-1-13 11](#_Toc449973259)

[2.1.28 TC-RNG-ATTR-2-13 11](#_Toc449973260)

[2.1.29 TC-STREAM-HASH-1-13 12](#_Toc449973261)

[2.1.30 TC-STREAM-HASH-2-13 12](#_Toc449973262)

[2.1.31 TC-STREAM-HASH-3-13 12](#_Toc449973263)

[2.1.32 TC-STREAM-ENC-1-13 12](#_Toc449973264)

[2.1.33 TC-STREAM-ENC-2-13 12](#_Toc449973265)

[2.1.34 TC-STREAM-ENCDEC-1-13 13](#_Toc449973266)

[3 KMIP Test Cases Setup 14](#_Toc449973267)

[3.1 KMIP 1.3 Test Cases Setup 14](#_Toc449973268)

[3.1.1 TC-CREG-1-13 14](#_Toc449973269)

[3.1.2 TC-CREG-3-13 14](#_Toc449973270)

[3.1.3 TC-OTP-SETUP-1-13 14](#_Toc449973271)

[3.1.4 TC-OTP-SETUP-2-13 15](#_Toc449973272)

[3.1.5 TC-OTP-CLEANUP-1-13 15](#_Toc449973273)

[Appendix A. Acknowledgments 16](#_Toc449973274)

[Appendix B. Revision History 18](#_Toc449973275)

1. Introduction

The purpose of this document is to describe test cases to demonstrate the Key Management Interoperability Protocol (KMIP) [KMIP-SPEC]. The test cases illustrate that the concepts within the protocol are sound and how the protocol may be used when implementing KMIP in applications. These test cases are not intended to fully test an implementation of KMIP.

## References (non-normative)

[KMIP-SPEC]

*Key Management Interoperability Protocol Specification Version 1.3*. Edited by Kiran Thota and Tony Cox. Latest version: <http://docs.oasis-open.org/kmip/spec/v1.3/kmip-spec-v1.3.html>.

[KMIP-PROFILES]

*Key Management Interoperability Protocol Profiles Version 1.3.* Edited by Tim Hudson and Robert Lockhart. Latest version: <http://docs.oasis-open.org/kmip/profiles/v1.3/kmip-profiles-v1.3.html>.

[XML]

*XML 1.0 Recommendation,* T. Bray, J. Paoli, M. Sperberg-McQueen, Editors, W3C Recommendation, February 10, 1998, <http://www.w3.org/TR/1998/REC-xml-19980210>. Latest version available at <http://www.w3.org/TR/REC-xml>.

1. KMIP Test Cases

The test cases define a number of request-response pairs for KMIP operations. Each test case is provided in the XML format specified in [KMIP-PROFILES] intended to be both human-readable and usable by automated tools.

Each test case has a unique label (the section name) which the protocol version as part of the identifier.

The test cases may depend on a specific configuration of a KMIP client and server being configured in a manner consistent with the test case assumptions.

Where possible the flow of unique identifiers between tests, the date-time values, and other dynamic items are indicated using symbolic identifiers – in actual request and response messages these dynamic values will be filled in with valid values.

The test cases show one possible way to construct the messages, and the messages shown are not necessarily the only conformant constructions as many items within KMIP are optional and server behavior depends on the server's policy. Support for a test case is predicated on a server matching the test case assumptions and the behavior shown in the request-response pairs.

Symbolic identifiers are of the form $UPPERCASE\_NAME followed by optional unique index value. Wherever a symbolic identifier occurs in a test cases the implementation must replace it with a reasonable appearing datum of the expected type. Time values can be specified in terms of an offset from the current time in seconds of the form $NOW or $NOW-n or $NOW+n.

## KMIP 1.3 Test Cases

### TC-CREG-2-13

Assuming that a KMIP server has set up a keypair and corresponding certificate (or will generate these on-the-fly) for a given one time credential (username and password or OTP value) return in a single request the public key, private key, and corresponding certificate for use in subsequent connections.

How the server creates the keypair and certificate is outside of the scope of KMIP (it MAY be performed via KMIP operations or via an entirely separate non-KMIP approach).

It is assumed that the server implements an appropriate policy to only accept the client provided credential once with a time limit on how soon the credential remains valid and that the public key, private key, and certificate will only be returned once. The server may elect to keep, archive, or destroy the managed objects after the client has completed this request.

See <test-cases/kmip-v1.3/TC-CREG-2-13.xml>

### TC-OFFSET-1-13

A client requests the server creates a number of symmetric keys and then uses the Offset parameter in Locate to return various items.

See <test-cases/kmip-v1.3/TC-OFFSET-1-13.xml>

### TC-OFFSET-2-13

A client requests the server creates a number of symmetric keys and then uses the Offset parameter in Locate to return various items.

See <test-cases/kmip-v1.3/TC-OFFSET-2-13.xml>

### TC-OTP-1-13

One-Time-Pad encryption - assuming pad has been setup

How the server sets up and operates the one time pad is outside of the scope of KMIP - this is just an example usage for testing the encrypt/decrypt mechanism.

A KMIP server can implement handling of the one-time-pad material via whatever approach makes sense in the context of a specific server implementation - all that is required is that both servers involved are in agreement about the one-time-pad.

See <test-cases/kmip-v1.3/TC-OTP-1-13.xml>

### TC-OTP-2-13

One-Time-Pad decryption - assuming pad has been setup

How the server sets up and operates the one time pad is outside of the scope of KMIP - this is just an example usage for testing the encrypt/decrypt mechanism.

A KMIP server can implement handling of the one-time-pad material via whatever approach makes sense in the context of a specific server implementation - all that is required is that both servers involved are in agreement about the one-time-pad.

See <test-cases/kmip-v1.3/TC-OTP-2-13.xml>

### TC-OTP-3-13

One-Time-Pad attempted get - assuming pad has been setup

Note: this example shows a server configured to return a Get without the key material present.

See <test-cases/kmip-v1.3/TC-OTP-3-13.xml>

### TC-OTP-4-13

One-Time-Pad attempted get - assuming pad has been setup

Note: this example shows a server configured to return denied for a Get request; the key material is never returned to the client in this configuration.

See <test-cases/kmip-v1.3/TC-OTP-4-13.xml>

### TC-OTP-5-13

One-Time-Pad attempted get - assuming pad has been setup and supports multiple encrypt and decrypt operations.

See <test-cases/kmip-v1.3/TC-OTP-5-13.xml>

### TC-Q-CAP-1-13

Return a list of responses indicating the server does not want to provide details as to its specific capabilities.

See <test-cases/kmip-v1.3/TC-Q-CAP-1-13.xml>

### TC-Q-CAP-2-13

Return a list of responses indicating the server does not want to provide details as to its specific capabilities.

See <test-cases/kmip-v1.3/TC-Q-CAP-2-13.xml>

### TC-Q-CREG-1-13

Return the list of client registration methods supported by a server. This example shows all four approaches are supported.

See <test-cases/kmip-v1.3/TC-Q-CREG-1-13.xml>

### TC-Q-PROF-1-13

Return details of the server claimed supported profiles.

See <test-cases/kmip-v1.3/TC-Q-PROF-1-13.xml>

### TC-Q-PROF-2-13

Return details of the server claimed supported profiles. This example shows a server claiming to support all profiles.

See <test-cases/kmip-v1.3/TC-Q-PROF-2-13.xml>

### TC-Q-PROF-3-13

 Return details of the server claimed supported profiles. This example shows a server returning Server URI and Port values for HTTPS usage

See <test-cases/kmip-v1.3/TC-Q-PROF-3-13.xml>

### TC-Q-RNGS-1-13

Return details of the supported RNGs where the server provides no actual information about the RNG (i.e. nothing is claimed).

See <test-cases/kmip-v1.3/TC-Q-RNGS-1-13.xml>

### TC-Q-RNGS-2-13

Return details of the supported RNGs where the server provides details of an ANSI X9.31 AES-256 based RNG. (e.g. RNGVAL 1202)

See <test-cases/kmip-v1.3/TC-Q-RNGS-2-13.xml>

### TC-Q-RNGS-3-13

Return details of the supported RNGs where the server provides details of an FIPS 186-2 x-Chagne Notice SHA-1 based RNG. (e.g. RNGVAL 1203)

See <test-cases/kmip-v1.3/TC-Q-RNGS-3-13.xml>

### TC-Q-RNGS-4-13

Return details of the supported RNGs where the server provides details of a DRBG HMAC based HMAC-SHA256 with prediction resistance RNG and a DRBG HMAC based HMAC-SHA1 with prediction resistance RNG and a DRBG Hash based SHA256 with prediction resistance RNG. (e.g. DRBGVAL 540)

See <test-cases/kmip-v1.3/TC-Q-RNGS-4-13.xml>

### TC-Q-RNGS-5-13

Return details of the supported RNGs where the server provides details of a DRBG Dual-EC based SHA-256 P-256 with prediction resistance RNG. (e.g. DRBGVAL 480)

See <test-cases/kmip-v1.3/TC-Q-RNGS-5-13.xml>

### TC-Q-RNGS-6-13

Return details of the supported RNGs where the server provides details of use of a plain AES-based DRBG

See <test-cases/kmip-v1.3/TC-Q-RNGS-6-13.xml>

### TC-Q-S2C-1-13

Server to Client Server queries the client's capabilities Client returns what it supports and may elect to use on the client to server link. This example is for a client supporting only the required operations and object types in the Tape Library Profile.

See <test-cases/kmip-v1.3/TC-Q-S2C-1-13.xml>

### TC-Q-S2C-2-13

Server to Client Server queries what KMIP protocol versions it supports Client returns the protocol versions it may use on the client to server link.

See <test-cases/kmip-v1.3/TC-Q-S2C-2-13.xml>

### TC-Q-S2C-PROF-1-13

Return details of the client claimed supported profiles. This is server-to-client request. Client returns the profiles it may use on the client to server link.

See <test-cases/kmip-v1.3/TC-Q-S2C-PROF-1-13.xml>

### TC-Q-S2C-PROF-2-13

Return details of the client claimed supported profiles. This is server-to-client request. Client returns the profiles it may use on the client to server link.

See <test-cases/kmip-v1.3/TC-Q-S2C-PROF-2-13.xml>

### TC-Q-VAL-1-13

Return details of the server claimed validation information. Example is for NIST CMVP FIPS140-2

See <test-cases/kmip-v1.3/TC-Q-VAL-1-13.xml>

### TC-Q-VAL-2-13

Return details of the server that does not claim any validations

See <test-cases/kmip-v1.3/TC-Q-VAL-2-13.xml>

### TC-RNG-ATTR-1-13

A client registers a symmetric key including details of the RNG that the client is claiming was used to generate the symmetric key.

See <test-cases/kmip-v1.3/TC-Q-RNG-ATTR-1-13.xml>

### TC-RNG-ATTR-2-13

A client requests the server creates a symmetric key and it does and also includes the required details of the RNG that was used to generate the symmetric key.

See <test-cases/kmip-v1.3/TC-Q-RNG-ATTR-2-13.xml>

### TC-STREAM-HASH-1-13

Hash operation for data 'abc' in a single request followed immediately by a streaming equivalent for which the result must be identical.

Note: - test vector data from http://csrc.nist.gov/groups/ST/toolkit/documents/Examples/SHA\_All.pdf

See <test-cases/kmip-v1.3/TC-STREAM-HASH-1-13.xml>

### TC-STREAM-HASH-2-13

Hash operation for data 'abc' in a single request followed immediately by a streaming equivalent for which the result must be identical.

Note: - test vector data from http://csrc.nist.gov/groups/ST/toolkit/documents/Examples/SHA\_All.pdf

See <test-cases/kmip-v1.3/TC-STREAM-HASH-2-13.xml>

### TC-STREAM-HASH-3-13

Hash operation for data 'abc' in a single request followed immediately by a streaming equivalent for which the result must be identical.

Note: - test vector data from http://csrc.nist.gov/groups/ST/toolkit/documents/Examples/SHA\_All.pdf

See <test-cases/kmip-v1.3/TC-STREAM-HASH-3-13.xml>

### TC-STREAM-ENC-1-13

Create a symmetric key and perform encrypt using the symmetric key.

Variation on CS-BC-M-1-13 in [KMIP-PROFILES]

See <test-cases/kmip-v1.3/TC-STREAM-ENC-1-13.xml>

### TC-STREAM-ENC-2-13

Create a symmetric key and perform encrypt using the symmetric key. Cryptographic Parameters are provided in first Encrypt operation rather than specified against the managed object.

Variation on CS-BC-M-4-13 in [KMIP-PROFILES]

See <test-cases/kmip-v1.3/TC-STREAM-ENC-2-13.xml>

### TC-STREAM-ENCDEC-1-13

Register a symmetric key and perform encrypt using the symmetric key followed by decrypt. The input data is non-block size.

Variation on CS-BC-M-10-13 in [KMIP-PROFILES]

See <test-cases/kmip-v1.3/TC-STREAM-ENC-3-13.xml>

1. KMIP Test Cases Setup

The test cases defined in the previous section all operate independent and assume that the other end of the KMIP connection has been configured to match the assumptions in the test case.

The following scripts allow for setting up the pre-conditions for a number of the test cases and for cleaning up after the test cases have executed – via KMIP operations. A server is not required to use KMIP or to use these scripts for this purpose – they are provided simply because they are useful for some implementations.

## KMIP 1.3 Test Cases Setup

### TC-CREG-1-13

This is used to set up the test data used in the client registration example. How the server sets up this in a normal context is outside of the scope of KMIP - this is just an example usage with configuration via KMIP with a pre-generated keypair and corresponding certificate.

A KMIP server can implement the equivalent capability via whatever approach makes sense in the context of a specific server implementation.

Register a public/private key pair in the PKCS\_1 key format and a corresponding X509 certificate. Add the appropriate links between the registered objects.

See <test-cases/kmip-v1.3/TC-CREG-1-13.xml>

### TC-CREG-3-13

This is used to clean up the test data used in the client registration example. How the server sets up this in a normal context is outside of the scope of KMIP - this is just an example usage with configuration via KMIP with a pre-generated keypair and corresponding certificate.

A KMIP server can implement the equivalent capability via whatever approach makes sense in the context of a specific server implementation.

See <test-cases/kmip-v1.3/TC-CREG-3-13.xml>

### TC-OTP-SETUP-1-13

One-Time-Pad setup - for testing purposes only

How the server sets up and operates the one time pad is outside of the scope of KMIP - this is just an example usage for testing the encrypt/decrypt mechanism where the one-time-pad has been set up with a simple value for testing rather than actually hooked into a real secure one-time-pad.

A KMIP server can implement handling of the one-time-pad material via whatever approach makes sense in the context of a specific server implementation - all that is required is that both servers involved are in agreement about the one-time-pad.

See <test-cases/kmip-v1.3/TC-OTP-SETUP-1-13.xml>

### TC-OTP-SETUP-2-13

One-Time-Pad setup - for testing purposes only

How the server sets up and operates the one time pad is outside of the scope of KMIP - this is just an example usage for testing the encrypt/decrypt mechanism where the one-time-pad has been set up with a simple value for tetsing rather than actually hooked into a real secure one-time-pad.

A KMIP server can implement handling of the one-time-pad material via whatever approach makes sense in the context of a specific server implementation - all that is required is that both servers involved are in agreement about the one-time-pad.

This setup allows for specification of the details without provision of the key material.

See <test-cases/kmip-v1.3/TC-OTP-SETUP-2-13.xml>

### TC-OTP-CLEANUP-1-13

One-Time-Pad cleanup - for testing purposes only

How the server sets up and operates the one time pad is outside of the scope of KMIP - this is just an example usage for testing the encrypt/decrypt mechanism where the one-time-pad has been set up with a simple value for tetsing rather than actually hooked into a real secure one-time-pad.

A KMIP server can implement handling of the one-time-pad material via whatever approach makes sense in the context of a specific server implementation - all that is required is that both servers involved are in agreement about the one-time-pad.

See <test-cases/kmip-v1.3/TC-OTP-CLEANUP-1-13.xml>

1. Acknowledgments

The following individuals have participated in the creation of this specification and are gratefully acknowledged:

**Editors of the previous versions of this document:**

Mathias Björkqvist, [IBM](http://www.ibm.com) (v1.0 and v1.1)

Tim Hudson, [Cryptsoft](http://www.cryptsoft.com) (v1.1)

René Pawlitzek, [IBM](http://www.ibm.com) (v1.0)

**Technical Committee Participants:**

Warren Armstrong, QuintessenceLabs Pty Ltd.

Rinkesh Bansal, IBM

Lina Baquero, Fornetix

Jeff Bartell, Fornetix

Tom Benjamin, IBM

Anthony Berglas, Cryptsoft Pty Ltd.

Mathias Bjorkqvist, IBM

Todd Bottger, Oracle

Joseph Brand, Semper Fortis Solutions

Alan Brown, Thales e-Security

Robert Burns, Thales e-Security

Andrew Byrne, EMC

Hai-May Chao, Oracle

Chye-Lin Chee, Hewlett Packard Enterprise (HPE)

Tim Chevalier, NetApp

Kenli Chong, QuintessenceLabs Pty Ltd.

Justin Corlett, Cryptsoft Pty Ltd.

Tony Cox, Cryptsoft Pty Ltd.

DINESH DIALANI, SafeNet, Inc.

Michael Dong, Hewlett Packard Enterprise (HPE)

Alex Downey, Futurex

Kevin Driver, IBM

Stephen Edwards, Fornetix

James Espinoza, Futurex

Faisal Faruqui, Thales e-Security

Stan Feather, Hewlett Packard Enterprise (HPE)

Indra Fitzgerald, NetApp

Judith Furlong, EMC

Michael Gardiner, SafeNet, Inc.

Jonathan Geater, Thales e-Security

Susan Gleeson, Oracle

Saheem Granados, IBM

John Green, QuintessenceLabs Pty Ltd.

Robert Griffin, EMC

Robert Haas, IBM

Steve He, Vormetric, Inc.

Christopher Hiller, Hewlett Packard Enterprise (HPE)

Hao Hoang, Hewlett Packard Enterprise (HPE)

Tim Hudson, Cryptsoft Pty Ltd.

Michael Jenkins, National Security Agency

Elysa Jones, Individual

Mark Joseph, P6R, Inc

Mahadev Karadigudda, NetApp

Jason Katonica, IBM

Tim Kelsey, Hewlett Packard Enterprise (HPE)

Hyun-jin Kim, Hancom Secure, Inc.

Stephen Kingston, SafeNet, Inc.

Kathy Kriese, Symantec Corp.

Leonardo Ladeira, SafeNet, Inc.

Sun-ho Lee, Hancom Secure, Inc.

John Leiseboer, QuintessenceLabs Pty Ltd.

Hal Lockhart, Oracle

Robert Lockhart, Thales e-Security

Martin Luther, Hewlett Packard Enterprise (HPE)

Jane Melia, QuintessenceLabs Pty Ltd.

Prashant Mestri, IBM

Trisha Paine, SafeNet, Inc.

Incheon Park, Hancom Secure, Inc.

John Peck, IBM

Michael Phillips, Dell

Stefan Pingel, EMC

Ajai Puri, SafeNet, Inc.

Saravanan Ramalingam, Thales e-Security

Bruce Rich, Cryptsoft Pty Ltd.

Warren Robbins, Dell

Peter Robinson, EMC

Rick Robinson, IBM

Saikat Saha, Oracle

Boris Schumperli, Cryptomathic

Greg Scott, Cryptsoft Pty Ltd.

Amit Sinha, SafeNet, Inc.

Radhika Siravara, Oracle

Curtis Smith, Futurex

Ryan Smith, Futurex

Amruta Soman, Cryptsoft Pty Ltd.

Gerald Stueve, Fornetix

Jim Susoy, P6R, Inc

Kiran Thota, VMware, Inc.

Peter Tsai, Vormetric, Inc.

Nathan Turajski, Hewlett Packard Enterprise (HPE)

Charles White, Fornetix

Steve Wierenga, Hewlett Packard Enterprise (HPE)

Thomas Xuelin, Watchdata Technologies Pte Ltd.

Krishna Yellepeddy, IBM

Magda, Zdunkiewicz, Cryptsoft Pty Ltd.

yuan zhang, Watchdata Technologies Pte Ltd.

Joshua Zhu, Vormetric, Inc.

1. Revision History

|  |  |  |  |
| --- | --- | --- | --- |
| **Revision** | **Date** | **Editor** | **Changes Made** |
| wd03 | 8-Jan-2016 | Tim Hudson  | Reformatted to have test cases in external files. |
| wd02 | 26-Jan-2015 | Tim Hudson / Mark Joseph | Updates to move setup related items in to a separate section. |
| wd01 | 22-Jan-2015 | Tim Hudson / Mark Joseph | Initial draft. |